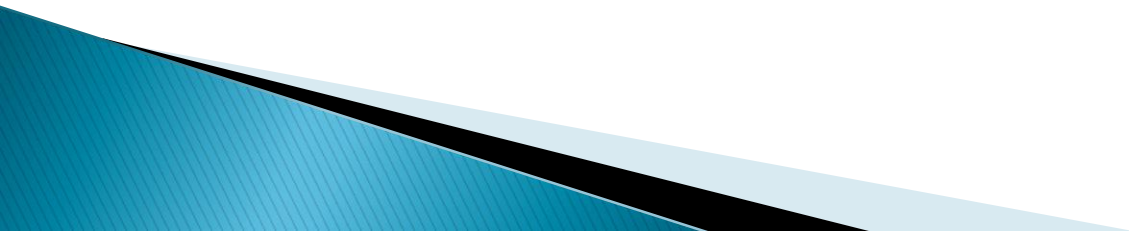


# Network Model and Hierarchy Model



# Introduction

- ▶ In network model the data is stored in the form of graph In the hierarchy model data is stored in the form of Trees

# Database Models

Network Model

Hierarchical Model

Relational Model

Object/Relational Model

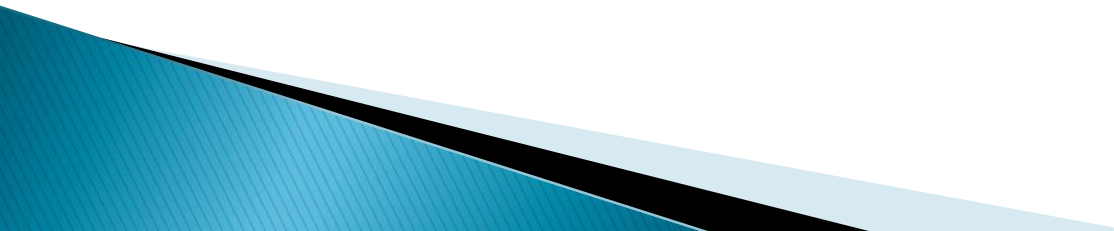
Object-Oriented Model

Semi structured Model

Associative Model

Entity-Attribute-Value (EAV) data model

Context Model



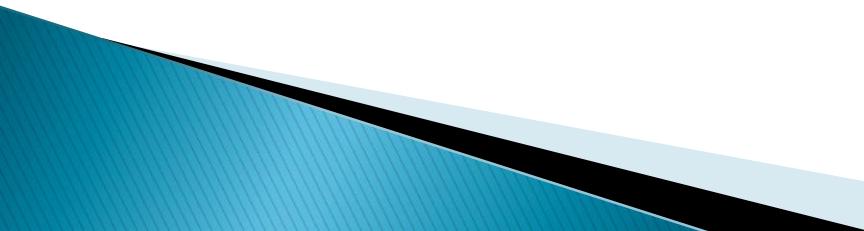
# Hierarchy Model

- ▶ The hierarchical data model organizes data in a tree structure. There is a hierarchy of parent and child data segments. This structure implies that a record can have repeating information, generally in the child data segments. Data in a series of records, which have a set of field values attached to it.

# continue

- ▶ It collects all the instances of a specific record together as a record type. These record types are the equivalent of tables in the relational model, and with the individual records being the equivalent of rows.

# Network Model

- ▶ Basic Concepts
  - ▶ Data-Structure Diagrams
  - ▶ The DBTG(database task group) CODASYL (multiuser compliance database mgmt system)Model
  - ▶ DBTG Data-Retrieval Facility
  - ▶ DBTG Update Facility
  - ▶ DBTG Set-Processing Facility
  - ▶ Mapping of Networks to Files
- 

# Basic Concepts

- ▶ Data are represented by collections of *records*.

- similar to an entity in the E-R model
- Records and their fields are represented as *record type*

```
type          customer = record  type    account =
record
  customer-name: string;          account-number: integer;
  customer-street: string;       balance: integer;
  customer-city: string;
end          end
```

- ▶ Relationships among data are represented by *links*
  - similar to a restricted (binary) form of an E-R relationship
  - restrictions on links depend on whether the relationship is many-many, many-to-one, or one-to-one.

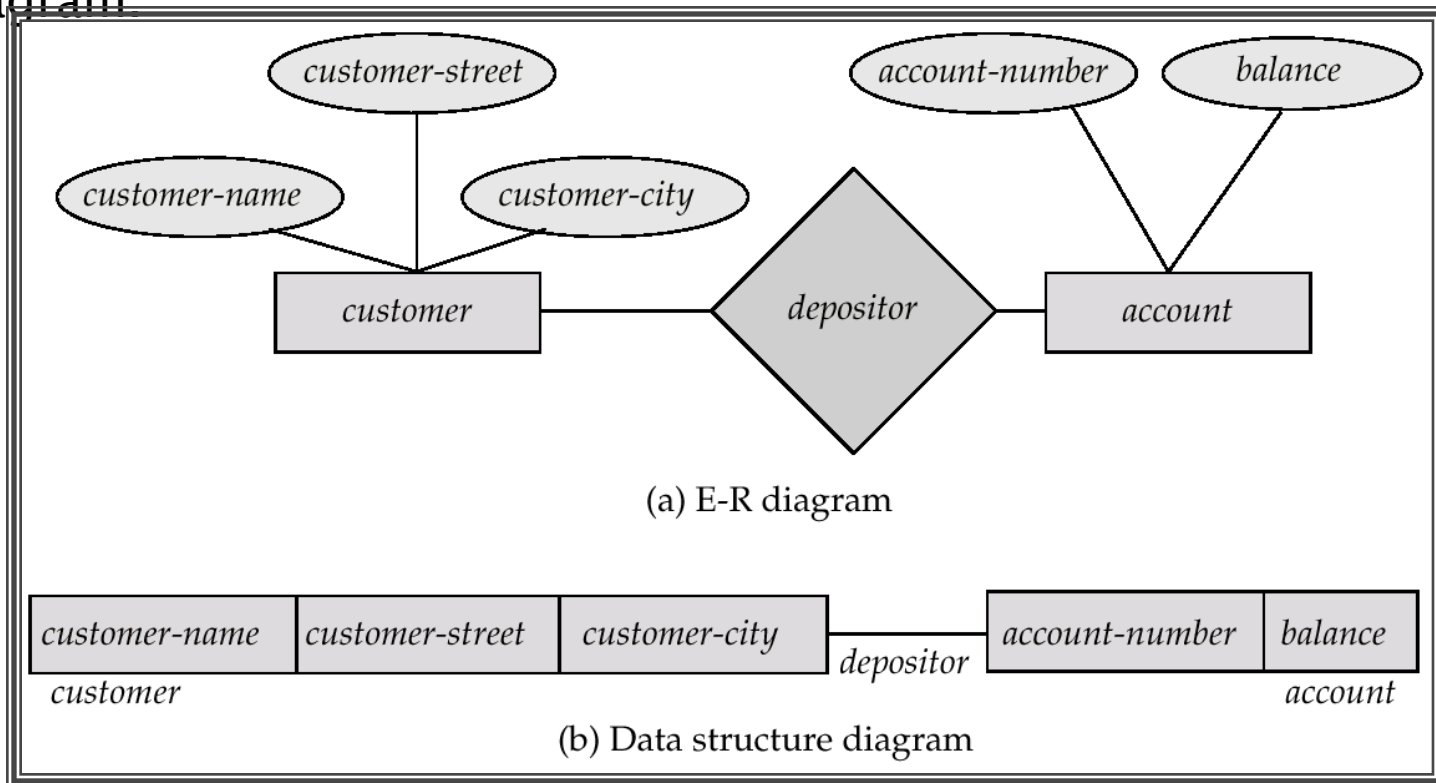
# Data-Structure Diagrams

- ▶ Schema representing the design of a network database.
- ▶ A data-structure diagram consists of two basic components:
  - **Boxes**, which correspond to record types.
  - **Lines**, which correspond to links.
- ▶ Specifies the overall logical structure of the database.



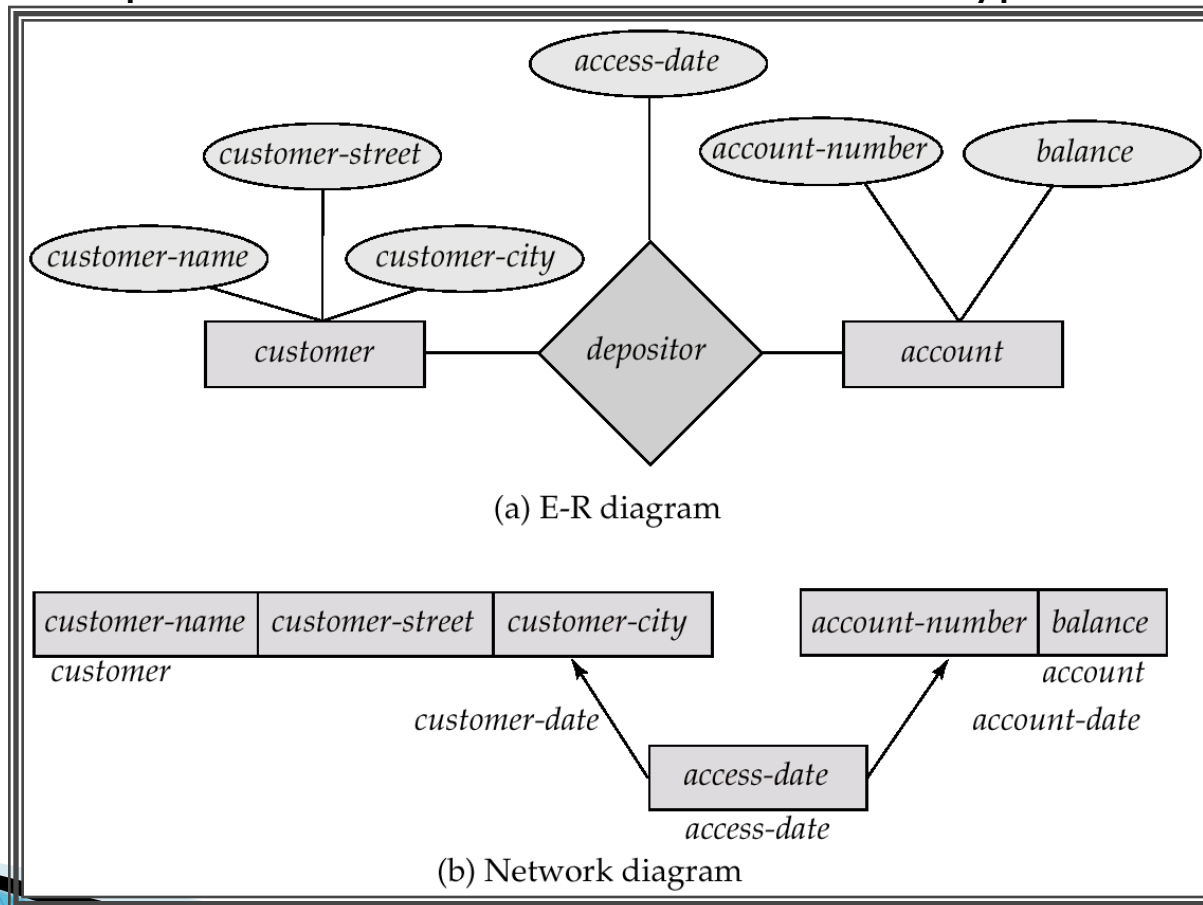
# Data-Structure Diagrams (Cont.)

- ▶ For every E-R diagram, there is a corresponding data-structure diagram



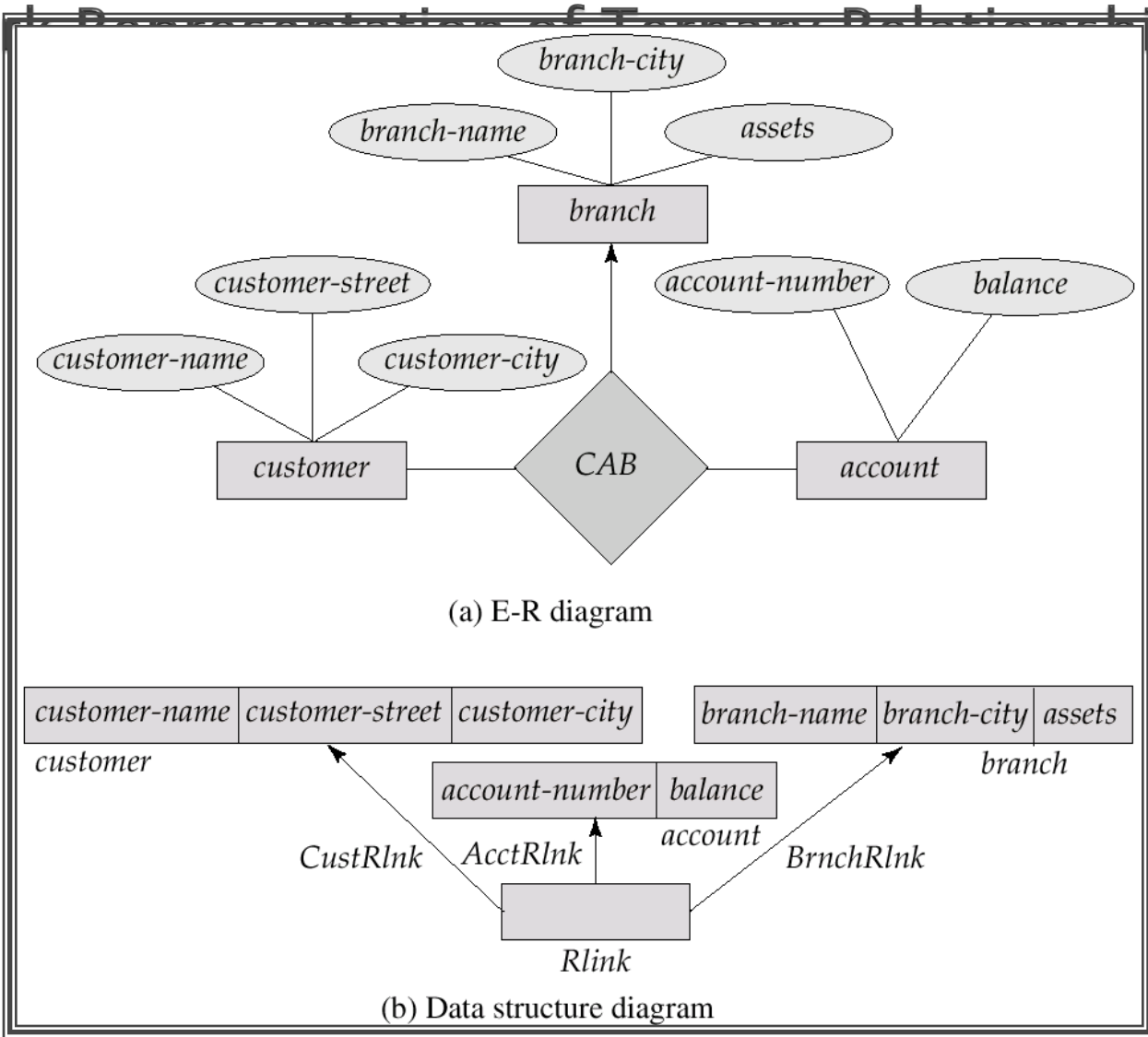
# Data-Structure Diagrams (Cont.)

- ▶ Since a link cannot contain any data value, represent an E-R relationship with attributes with a new record type and links.



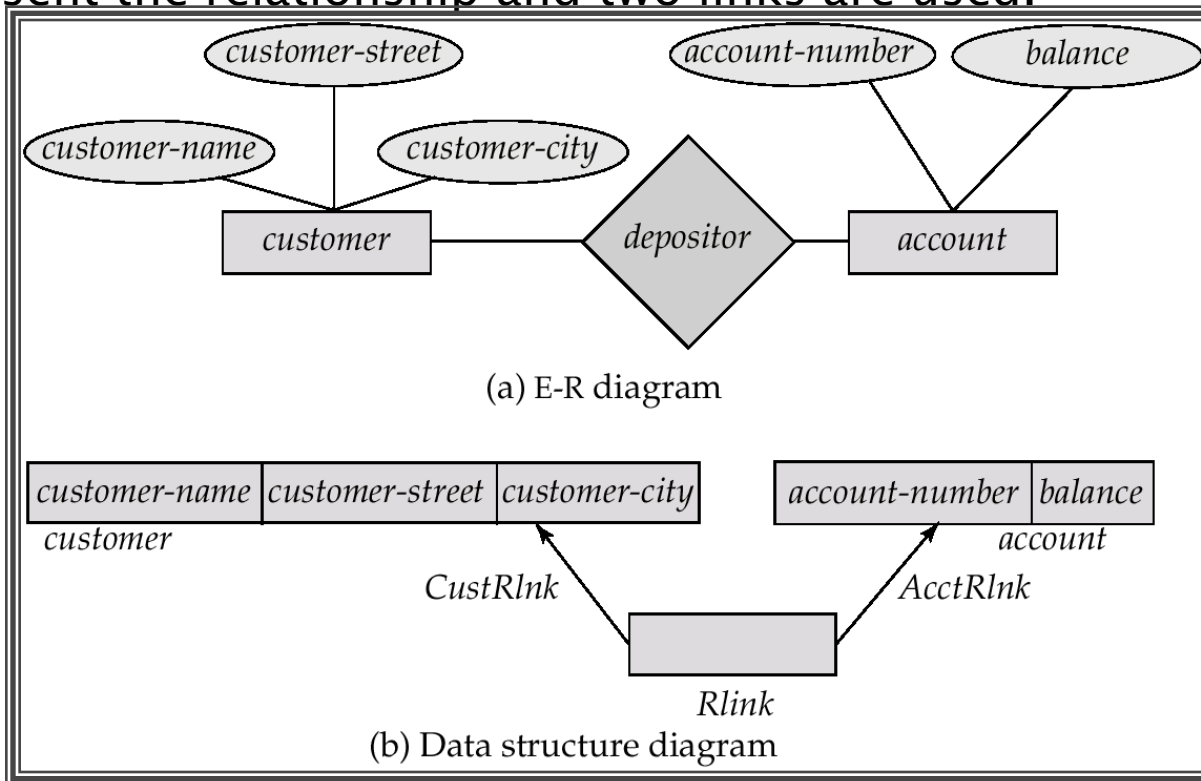
# General Relationships

- ▶ To represent an E–R relationship of degree 3 or higher, connect the participating record types through a new record type that is linked directly to each of the original record types.
  1. Replace entity sets *account*, *customer*, and *branch* with record types *account*, *customer*, and *branch*, respectively.
  2. Create a new record type *Rlink* (referred to as a *dummy* record type).
  3. Create the following many–to–one links:
    - *CustRlink* from *Rlink* record type to *customer* record type
    - *AcctRlink* from *Rlink* record type to *account* record type
    - *BrncRlink* from *Rlink* record type to *branch* record type

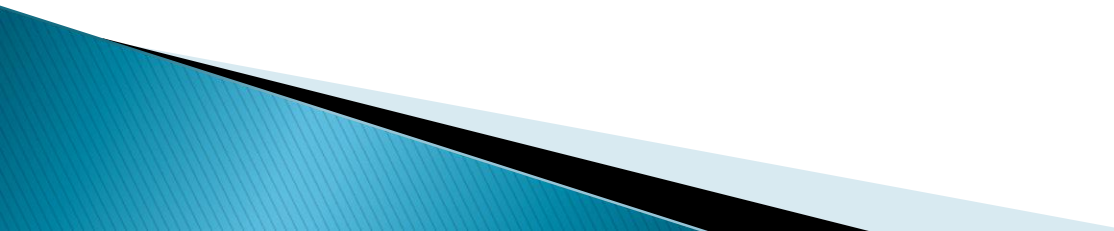


# The DBTG CODASYL Model

- ▶ All links are treated as many-to-one relationships.
- ▶ To model many-to-many relationships, a record type is defined to represent the relationship and two links are used.



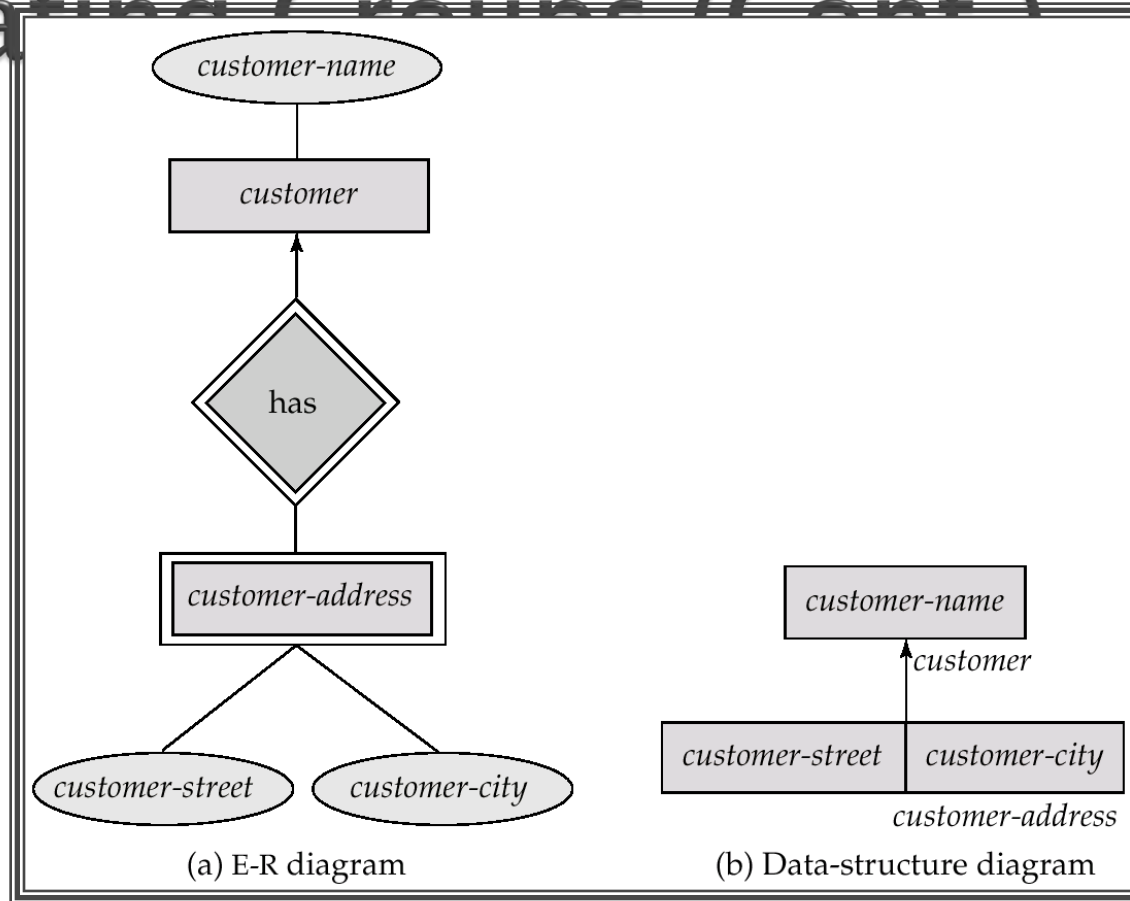
# DBTG Sets

- ▶ The structure consisting of two record types that are linked together is referred to in the DBTG(database task group) model as a *DBTG set*
  - ▶ In each DBTG set, one record type is designated as the *owner*, and the other is designated as the *member*, of the set.
  - ▶ Each DBTG set can have any number of *set occurrences* (actual instances of linked records).
  - ▶ Since many-to-many links are disallowed, each set occurrence has precisely one owner, and has zero or more member records.
  - ▶ No member record of a set can participate in more than one occurrence of the set at any point.
  - ▶ A member record can participate simultaneously in several set occurrences of *different* DBTG sets.
- 

# Repeating Groups

- ▶ Provide a mechanism for a field to have a set of values rather than a single value.
- ▶ Alternative representation of weak entities from the E-R model
- ▶ Example: Two sets.
  - *customer (customer-name)*
  - *customer-address (customer-street, customer-city)*
- ▶ The following diagrams represent these sets without the repeating-group construct.

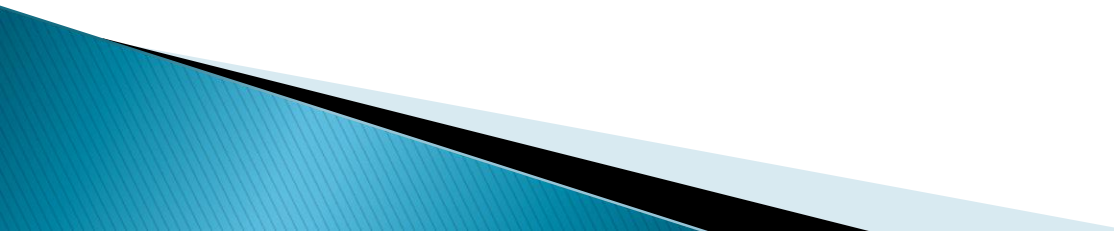
# Repeating Groups (Cont)



- ▶ With the repeating-group construct, the data-structure diagram consists of the single record type *customer*.



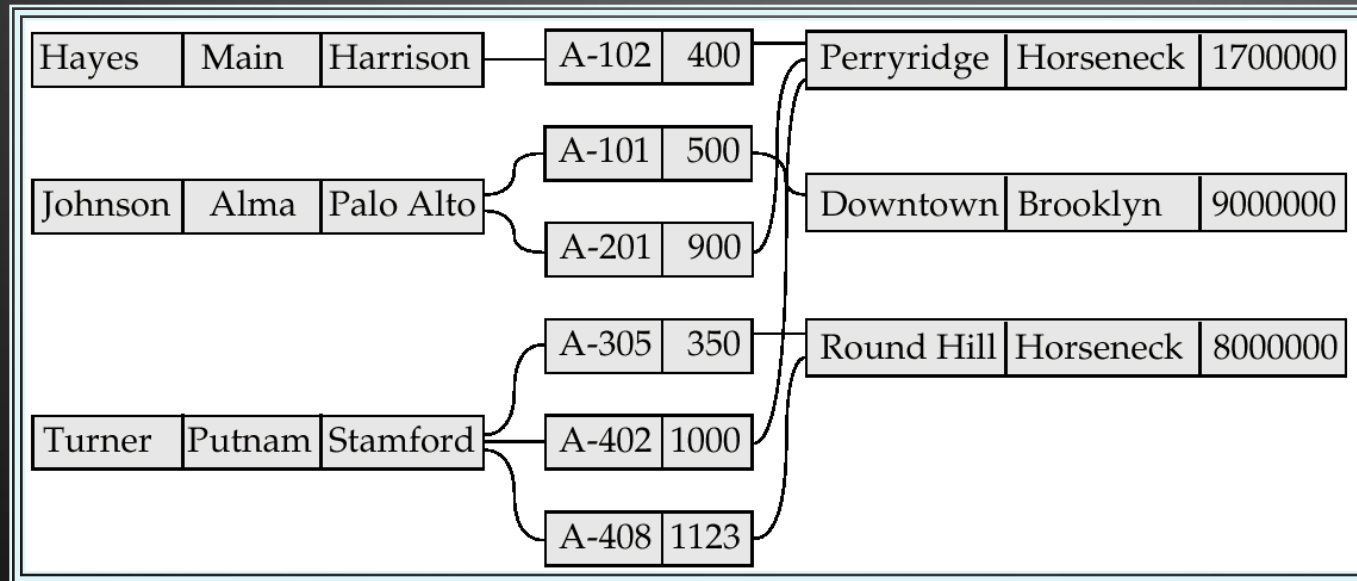
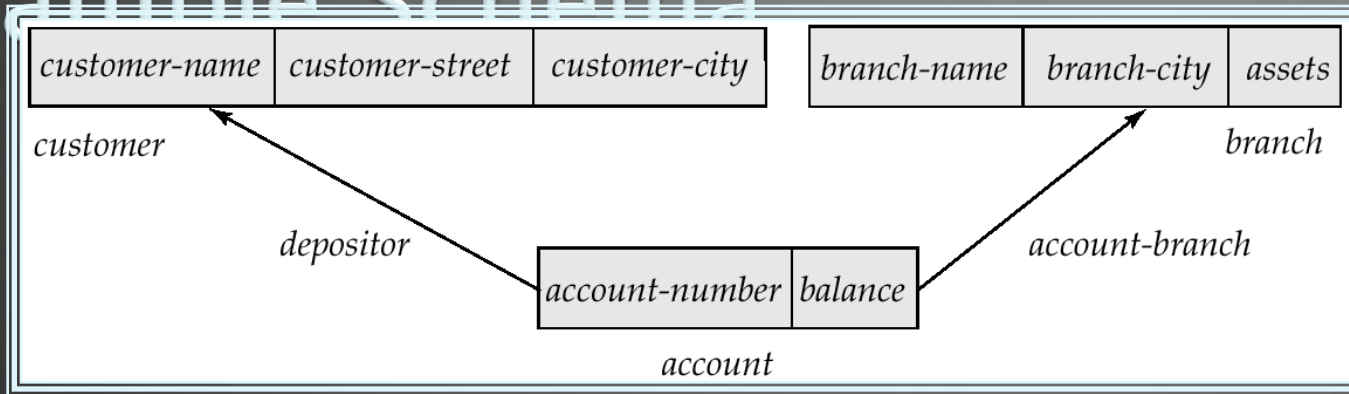
# DBTG Data–Retrieval Facility

- ▶ The DBTG data manipulation language consists of a number of commands that are embedded in a host language.
  - ▶ *Run unit* — system application program consisting of a sequence of host language and DBTG command statements. Statements access and manipulate database items as well as locally declared variables.
  - ▶ *Program work–area* (or *user work area*) — a buffer storage area the system maintains for each application program
- 

# DBTG Variables

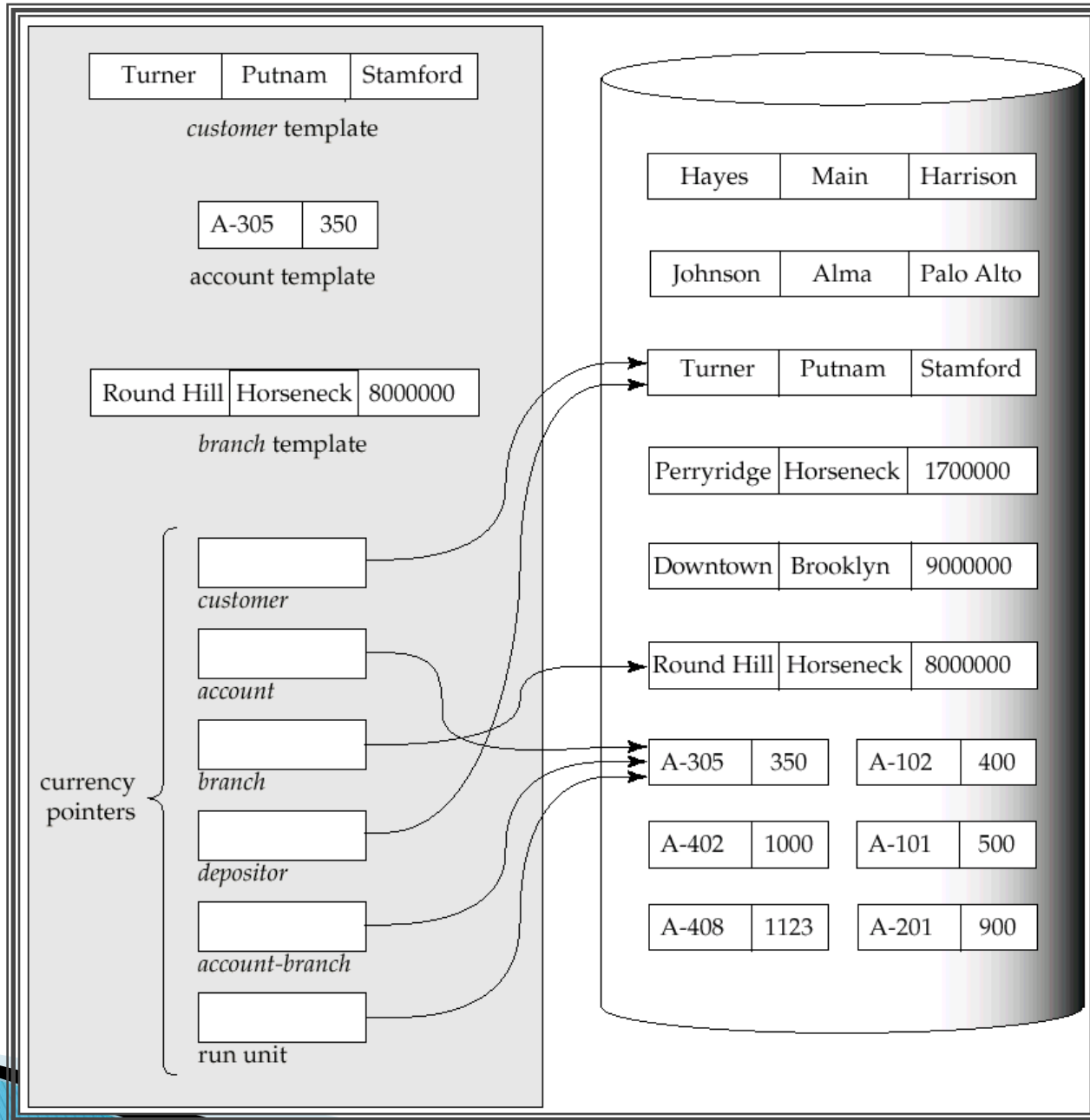
- ▶ Record Templates
- ▶ Currency pointers
  - Current of record type
  - Current of set type
  - Current of run unit
- ▶ Status flags
  - **DB-status** is most frequently used
  - Additional variables: **DB-set-name**, **DB-record-name**, and **DB-data-name**

# Example Schema



# Example Program Work Area

- ▶ Templates for three record types: *customer*, *account*, and *branch*.
- ▶ Six currency pointers
  - Three pointers for record types: one each to the most recently accessed *customer*, *account*, and *branch* record
  - Two pointers for set types: one to the most recently accessed record in an occurrence of the set *depositor*, one to the most recently accessed record in an occurrence of the set *account-branch*
  - One run-unit pointer.
- ▶ Status flags: four variables defined previously
- ▶ Following diagram shows an example program work area state.



# The Find and Get Commands

- ▶ **find** locates a record in the database and sets the appropriate currency pointers
- ▶ **get** copies of the record to which the current of run-unit points from the database to the appropriate program work area template
- ▶ Example: Executing a **find** command to locate the customer record belonging to Johnson causes the following changes to occur in the state of the program work area.
  - The current of the record type *customer* now points to the record of Johnson.
  - The current of set type *depositor* now points to the set owned by Johnson
  - The current of run unit now points to *customer* record Johnson.

# Access of Individual Records

- ▶ **find any** <record type> **using** <record-field>  
Locates a record of type <record type> whose <record-field> value is the same as the value of <record-field> in the <record type> template in the program work area.
- ▶ Once such a record is found, the following currency pointers are set to point to that record:
  - The current of run-unit pointer
  - The record-type currency pointer for <record type>
  - For each set in which that record belongs, the appropriate set currency pointer
- ▶ **find duplicate** <record type> **using** <record-field>  
Locates (according to a system-dependent ordering) the next record that matches the <record-field>

# Access of Records Within a Set

- ▶ Other **find** commands locate records in the DBTG set that is pointed to by the <set-type> currency pointer.
- ▶ **find first** <record type> **within** <set-type>  
Locates the first database record of type <record type> belonging to the current <set-type>.
- ▶ To locate the other members of a set, we use

**find next** <record type> **within** <set-type>

which finds the next element in the set <set-type>.

- ▶ **find owner within** <set-type>  
Locates the owner of a particular DBTG set



# Predicates

- ▶ For queries in which a field value must be matched with a specified range of values, rather than to only one, we need to:
  - **get** the appropriate records into memory
  - examine each one separately for a match
  - determine whether each is the; target of our **find** statement

# Example DBTG Query

- ▶ Print the total number of accounts in the Perryridge branch with a balance greater than \$10,000.

```
count := 0;  
branch.branch-name := "Perryridge";  
find any branch using branch-name;  
find first account within account-branch;  
while DB-status = 0 do  
  begin  
    get account  
    if account.balance > 10000 then count := count + 1;  
    find next account within account-branch;  
  end  
print (count);
```

# DBTG Update Facility

- ▶ DBTG mechanisms are available to update information in the database.
- ▶ To create a new record of type <record type>
  - insert the appropriate values in the corresponding <record type> template
  - add this new record to the database by executing  
**store <record type>**
- ▶ Can create and add new records only one at a time

# DBTG Update Facility (Cont.)

- ▶ To modify an existing record of type <record type>
  - find that record in the database
  - get that record into memory
  - change the desired fields in the template of <record type>
  - reflect the changes to the record to which the currency point of <record type> points by executing

**modify** <record type>

# DBTG Update Facility (Cont.)

- ▶ To delete an existing record of type <record type>
  - make the currency pointer of that type point to the record in the database to be deleted
  - delete that record by executing

**erase** <record type>

- ▶ Delete an entire set occurrence by finding the owner of the set and executing

**erase all** <record type>

- Deletes the owner of the set, as well as all the set's members.
- If a member of the set is an owner of another set, the members of that second set also will be deleted.
- **erase all** is recursive.

# DBTG Set-Processing Facility

- ▶ Mechanisms are provided for inserting records into and removing records from a particular set occurrence
- ▶ Insert a new record into a set by executing the **connect** statement.  
**connect** <record type> **to** <set-type>
- ▶ Remove a record from a set by executing the **disconnect** statement.  
**disconnect** <record type> **from** <set-type>